

The Maximal Defect k -Subgroups of Semigroups of Graphs and Digraphs

Károly Podoski



3rd BIOMICS Workshop
Passau, 9 February 2016

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Graphs and Digraphs

- ▶ *Digraph*: $\Gamma = (V, E)$, with vertices V , edges $E \subseteq V \times V$.
- ▶ The *reverse edge* to $e = (x, y) \in E$ is $\bar{e} = (y, x)$
- ▶ *(Undirected) Graph*: E is symmetric
- ▶ No self-loops: $(x, x) \notin E$.

Spectrum of Graph

Algebraic graph theory

- ▶ *adjacency matrix* $A = A(\Gamma)$ of $\Gamma = (V, E)$ is a matrix of 0's and 1's with $A_{x,y} = 1$ iff $(x, y) \in E$.
- ▶ The multiset of eigenvalues of A is called the *spectrum* of A .

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The spectrum of graphs and digraphs is invariant for isomorphic (di)graphs.

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$K_{1,4}$ and $C_4 \sqcup K_1$ have same spectrum. \implies **not** complete invariant

Elementary Collapsings of Directed Edges

Digraph $\Gamma = (V, E)$

Elementary collapsing

edge $e = (x, y) \longrightarrow$ function $T_{x,y}: V \rightarrow V$

$$T_{x,y}(v) = \begin{cases} y & \text{if } v = x, \\ v & \text{otherwise.} \end{cases}$$

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Example

$$T_{(1,2)}: 1 \rightarrow 2$$

$$2 \rightarrow 2$$

$$3 \rightarrow 3$$

$$4 \rightarrow 4$$

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$$T_{(1,2)}$$

Motivation: Biochemical Reactions

Biochemical transitions are modelled as products of commuting elementary collapsings, $f = \prod T_{a,b}$, where $T_{x,y}$ and $T_{y,z}$ do not both occur among the $T_{a,b}$ for any x, y , and z .

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We say a mapping $f : V \rightarrow V$ has *defect* k if $|f(V)| = |V| - k$.

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f has defect $k \implies fg$ and gf have defect at least k .

Composition of Elementary Collapsings

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Composition of Elementary Collapsings

Definition

The *transformation semigroup* of a digraph $\Gamma = (V, E)$ to be the collection of transformations of V generated by the T_e ($e \in E$).

$$S(\Gamma) = \langle T_{x,y} \in V^V \mid (x,y) \text{ is an edge of } \Gamma \rangle.$$

This is also called the **semigroup of flows** on Γ .

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Theorem

$S(\Gamma)$ is an invariant for graphs and for digraphs.

Proof.

Isomorphic (di)graphs obviously have isomorphic semigroups of flows. □

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Theorem (Nehaniv–Rhodes)

$S(\Gamma)$ is a complete algebraic invariant for undirected graphs.

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There exist non-isomorphic digraphs with the same flow semigroup.
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Now $f^2 = T_{3,2}$, corresponding to an edge (3,2).

Complete Invariance for Graphs

Lemma

If $T_{x,y} = T_{x_1,y_1} \cdots T_{x_k,y_k}$, then $T_{x,y}$ or $T_{y,x}$ appears among T_{x_i,y_i} .

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Corollary

Let Γ be a digraph.

If $T_{x,y} \in S(\Gamma)$ then either $e = (x,y)$ or $\bar{e} = (y,x)$ is an edge of Γ .

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Reversing Edges in Directed Cycles

Δ *directed n -cycle*: n nodes, edges $(1, 2), \dots, (n - 1, n), (n, 1)$.

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So, $S(\Delta \cup \{(1, n)\}) = S(\Delta)$.

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Reversing Edges in Directed Cycles

Theorem

Digraph Γ . If e is an edge in a simple directed cycle in Γ , then $S(\Gamma) = S(\Gamma \cup \{\bar{e}\})$. \implies Enough to consider undirected graphs on strongly connected components.

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$$i \rightarrow i+1 \quad \text{W defect 1: permutes } n-1 \text{ vertices cyclically}$$

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Defect k Permutator Groups

Permutator Group

- ▶ If $e^2 = e: V \rightarrow V$ is an idempotent $S(\Gamma)$.
Let $X_e = V \cdot e = \{v \cdot e \mid v \in V\}$
Let $G_e =$ (unique) maximal subgroup of $S(\Gamma)$ containing e .
- ▶ (X_e, G_e) (faithful) *permutator group* of subset X_e , consists of defect k where $k = |V| - |X_e|$.

Theorem

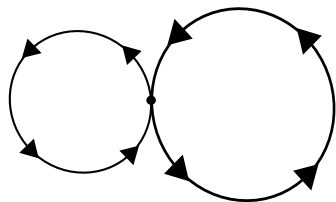
The defect k permutator groups (up to isomorphism) are invariants for digraphs.

Some Defect 1 Permutator Groups

Example

- ▶ Cycle graph with n nodes: Defect 1 group: cyclic C_{n-1}

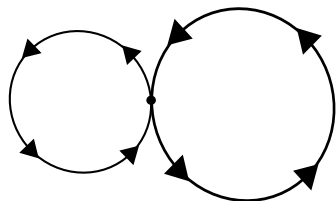
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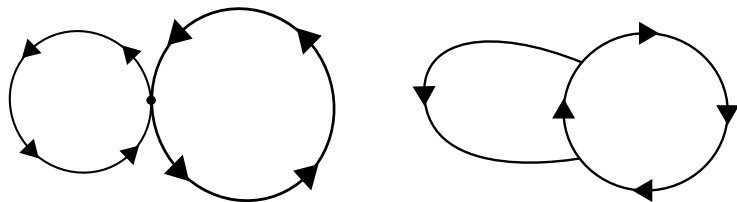
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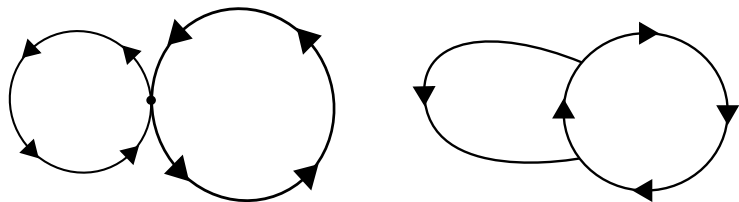
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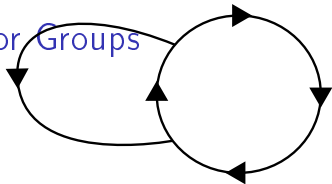


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- ▶ Two cycle graphs n and m nodes: Defect 1 group:
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- ▶ Overlapping Cycles: Defect 1 group: S_{n-1} or A_{n-1} or S_5

Defect 1 Permutator Groups

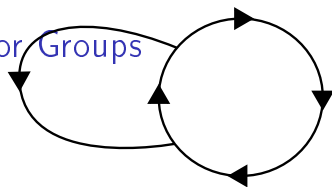
Theorem



$$\begin{aligned}T_{k,l,m} &= \langle x, y \rangle, n = k + l + m, \\x &= (a_1, a_2, \dots, a_k, b_1, b_2, \dots, b_l), \\y &= (a_1, a_2, \dots, a_k, c_1, c_2, \dots, c_m),\end{aligned}$$

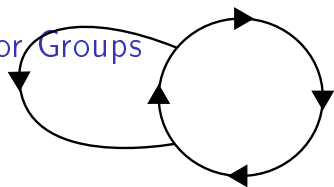
1. If $k = 0$ then $T_{0,l,m} \simeq C_l \times C_m$;
2. if $k \geq 1$, $k + l$ and $k + m$ are both odd, then $T_{k,l,m} = A_n$;
3. $T_{3,2,1} \simeq T_{2,2,2} \simeq T_{3,1,2} \simeq S_5$, and this is a 3-transitive action of S_5 on 6 elements;
4. $T_{k,l,m} = S_n$, otherwise.

Defect 1 Permutator Groups



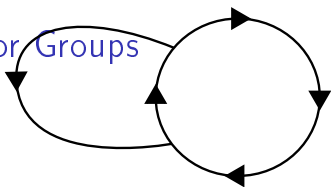
- ▶ ear decomposition

Defect 1 Permutator Groups



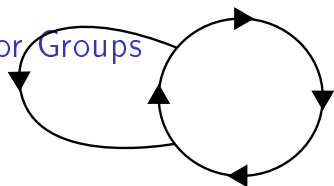
- ▶ ear decomposition \iff 2-edge connected graph

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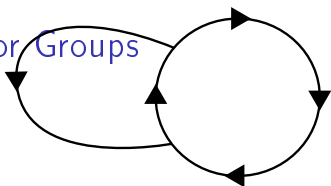
- ▶ ear decomposition \iff 2-edge connected graph
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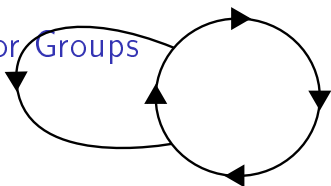
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- ▶ open ear decomposition

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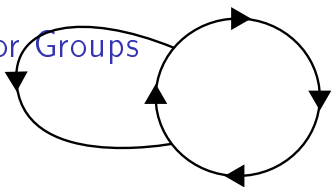
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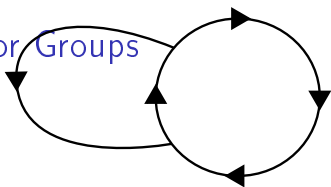
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- ▶ done for all 2-vertex connected graphs

Main Theorem (Horváth, Nehaniv, Podoski, 2015)

Graph Γ , 2-edge connected, not a cycle.

Then maximal defect k groups are all isomorphic and

- ▶ Defect 1: Direct product of cyclic, alternating and symmetric groups of various orders, (corresponding to the biconnected components)
- ▶ Defect $k \geq 2$, no bridge with $\geq k$ nodes $\implies S_{n-k}$
- ▶ Defect $k \geq 2$, bridges with $\geq k$ nodes \implies direct product of symmetric groups (corresponding to the biconnected components with $\leq k - 1$ long bridges)

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